

## IN THE CLAIMS

This listing of the claim will replace all prior versions and listings of claim in the present application.

### Listing of Claims

1. (currently amended) An information processing method for calculating  $x \cdot (2^n) \bmod P$  for an input value  $x$  larger than a prime number  $P$ , the operator  $\wedge$  denoting power, said information processing method being executed by an information processing apparatus comprising a first memory of  $n$  bits sufficient for storing the modulus  $P$ , a second memory of  $m$  bits sufficient for storing said input value  $x$ , a third memory for storing  $2^{(2m+n)} \bmod P$  and a Montgomery modular multiplication device, said information processing method comprising the steps ofwherein:

~~the value  $x \cdot (2^n) \bmod P$  is calculated without explicitly obtaining  $x \bmod P$ , by:~~ storing said input value  $x$  in said first memory of  $n$  bits;

~~calculating or previously preparing  $2^{(2m+n)} \bmod P$  or reading said~~ when the input value  $x$  has to be transformed into  $x \cdot (2^n) \cdot 2^{(2m+n)} \bmod P$  from said third memory by said Montgomery modular multiplication device;

~~reading said input value  $x$  from said first memory of  $n$  bits;~~

~~the number  $n$  denoting the number of bits necessary and sufficient for storing the modulus  $P$  and the number  $m$  denoting the number of bits necessary for storing the input value  $x$ ;~~

~~calculating  $x_1 = x \cdot 2^{(2m+n)} \cdot (2^{(-m)}) \bmod P = x \cdot 2^{(m+n)} \bmod P$  by~~ said Montgomery modular multiplication device; and

~~calculating  $x_2 = x_1 \cdot (2^{(-m)}) \bmod P = x \cdot (2^n) \bmod P$  to transfer said~~ input value  $x$  into  $x \cdot (2^n) \bmod P$  without explicitly obtaining  $x \bmod P$ ; and

storing said transferred value.

2. (currently amended) An information processing method for calculating  $x \cdot (2^n) \bmod P$  for an input value  $x$  larger than a prime number  $P$ , the operator  $^$  denoting power, said information processing method being executed by an information processing apparatus comprising a first memory of  $n$  bits sufficient for storing the modulus  $P$ , a second memory of  $m$  bits sufficient for storing said input value  $x$ , a third memory for storing  $2^{(m+2n)} \bmod P$  and a Montgomery modular multiplication device, said information processing method comprising the steps ofwherein: the value  $x \cdot (2^n) \bmod P$  is calculated without explicitly obtaining  $x \bmod P$ , by:

storing said input value  $x$  in said first memory of  $n$  bits;

calculating or previously preparing  $2^{(m+2n)} \bmod P$  or reading said  $2^{(m+2n)} \bmod P$  from said third memory by said Montgomery modular multiplication device;

reading said input value  $x$  from said first memory of  $n$  bits;  
~~when the input value  $x$  has to be transformed into  $x \cdot (2^n) \bmod P$ , the number  $n$  denoting the number of bits necessary and sufficient for storing the modulus  $P$  and the number  $m$  denoting the number of bits necessary for storing the input value  $x$ ;~~

calculating  $x_1 = x \cdot 2^{(m+2n)} \cdot (2^{(-m)}) \bmod P = x \cdot 2^{(2n)} \bmod P$  by said Montgomery modular multiplication device; and

calculating  $x_2 = x_1 \cdot (2^{(-n)}) \bmod P = x \cdot (2^n) \bmod P$  to transfer said input value  $x$  into  $x \cdot (2^n) \bmod P$  without explicitly obtaining  $x \bmod P$ ; and

storing said transferred value.

Claim 3 (canceled).

4. (currently amended) The information processing method of claim 1 for an RSA cryptosystem method using Chinese Remainder Theorem, said method comprising the steps of:

~~inputting an input value X;~~

calculating mod P using the information processing method according to claim 1 and encrypting said input value ~~the x~~; and

storing said ~~the~~ encrypted input value x.

5. (currently amended) The information processing method of claim 2 for an RSA cryptosystem method using Chinese Remainder Theorem, said method further comprising the steps of:

~~inputting an input value X;~~

calculating mod P using the information processing method according to claim 2, and encrypting said input value ~~the x~~; and

storing said ~~the~~ encrypted input value x.

6. (currently amended) An information processing apparatus for calculating  $x \cdot (2^n) \bmod P$  for an input value x larger than a prime number P, the operator ^ denoting power, said apparatus comprising:

a memory of n bits ~~wherein the number n denotes the number of bits necessary and sufficient for storing the modulus P;~~

a memory of m bits sufficient for ~~and the number m denotes the number of bits necessary for storing the~~ said input value x, and

a wherein the information processing apparatus comprises  
Montgomery modular multiplication device, wherein said and the Montgomery  
modular multiplication device adapted to execute the steps of  
calculating  $2^{(2m+n)} \bmod P$ ;  
reading said input value x from said memory of n bits; and  
calculating  $x_1 = x \cdot 2^{(2m+n)} \cdot (2^{-m}) \bmod P = x \cdot 2^{(m+n)} \bmod P$   
and  $x_2 = x_1 \cdot (2^{-m}) \bmod P = x \cdot (2^n) \bmod P$  to transfer said  
input value x into  $x \cdot (2^n) \bmod P$  without explicitly obtaining  $x \bmod P$ .